

# DATA SECURITY USING CRYPTOGRAPHY AND STEGANOGRAPHY TECHNIQUES

<sup>1</sup>P.V.N RAJESWARI, <sup>2</sup>SYED. KARISHMA SULTANA

<sup>1</sup>Assoc. Professor, Dept. of CSE, PBR VITS, Kavali, A.P, India.

<sup>2</sup>MCA, Dept. of MCA, PBR VITS, Kavali, A.P, India.

**Abstract** — Although cryptography and steganography could be used to provide data security, each of them has a problem. Cryptography problem is that, the cipher text looks meaningless, so the attacker will interrupt the transmission or make more careful checks on the data from the sender to the receiver. Steganography problem is that once the presence of hidden information is revealed or even suspected, the message is become known. According to the work in this paper, a merged technique for data security has been proposed using Cryptography and Steganography techniques to improve the security of the information. Firstly, the Advanced Encryption Standard (AES) algorithm has been modified and used to encrypt the secret message. Secondly, the encrypted message has been hidden using method in [1]. Therefore, two levels of security have been provided using the proposed hybrid technique. In addition, the proposed technique provides high embedding capacity and high quality stego images.

**Keywords**—Image Steganography; Pixel Value Difference (PVD); Encryption; Decryption; Advance encryption standard (AES)

## I. INTRODUCTION

Information transmission through internet may include sensitive personal data which may be intercepted. Also, there are many applications on the internet and many web sites require the users to fill forms that include sensitive personal information such as telephone numbers, addresses, and credit card information. So, the users may need private and secure communications for many reasons such as protect their confidential information from hackers during it passed over an open channel, so the confidentiality and data integrity are required to protect against unauthorized access and use. Cryptography and steganography are the common methods to secure communications [2]. Cryptography is the science of using mathematics to encrypt and decrypt data to keep messages secured by

transforming intelligible data form (plaintext) into unintelligible form (ciphertext). The term cryptography has come from the Greek word “kryptós” standing for “hidden” and “gráphin” standing for “writing”. Thus, the proper meaning of cryptography is “hidden writing” [3, 4]. Any cryptosystem consists of plaintext, encryption algorithm, decryption algorithm, Cipher text, and Key. Plaintext is message or data which are in their normal, readable (not encrypted) form. Encryption is the process of converting plaintext to cipher text by using key. Cipher text results from encryption by applying the encryption key on the plaintext. Decryption is the process of retrieving the plaintext back from the cipher text. The Key is used info to control the cryptosystem (cipher system), and it is known by the sender and receiver only [3, 5]. While cryptography is very powerful for securing data; the cryptanalysts could success to break the ciphers by analysing the contents of cipher text to get back the plaintext [3]. Cryptographic systems are generally classified into three independent dimensions [3]:

### A. Type of Operation on Plaintext

There are two types of operations that are occurred on plaintext to transform plaintext to cipher text. According to the first operation, each element in plaintext (i.e., bit, letter, group of bits or letters) is substituted for one another in the ciphertext. In this type of operation, a one-to-one mapping between the elements such as Caesar cipher [5]. The principle of the second type of operation is that each character in plaintext is transposed with one another based on a mapping dictated by the key. In this type, the plaintext characters stay the same but they are just moved into different positions such as Rail Fence cipher. Most systems, referred to as product systems, involve multiple stages of substitutions and transpositions.

*B. The Number of Used Keys*

If the sender and the receiver use one key to encrypt and decrypt the plaintext, the system is referred to as symmetric, single key, secret key or conventional encryption. Symmetric encryption is fairly straightforward and very fast. If the sender and receiver use different keys, public key and private key, to encrypt and decrypt the plaintext respectively, the system is referred to as asymmetric, two – key, or public key encryption.

*C. The Way in which The Plain Text is processed*

Block cipher operates on fixed-length groups of bits, called blocks, and produces an output block for each input block. A stream cipher operates on each plaintext element continuously, and produces one element at a time, as it goes along. On the other hand, Steganography is considered the art and science of hiding information in other information. The word Steganography is derived from the Greek words “steganos” meaning “impenetrable” and, “grafia” meaning “writing” defining it as “impenetrable writing” [4, 6]. There are two common techniques for image embedding in steganography; spatial domain and transform domain. According to spatial domain embedding, the messages are embedded directly into the Least Significant Bits (LSBs). The least significant bits (LSB) insertion method is considered the most common and simplest Steganography method. According to transform domain embedding, the messages are embedded by modifying frequency coefficients of the cover image such as the Fourier transform, discrete cosine transform, or the wavelet transform [7]. Image steganography system is comprised two algorithms, one for embedding and one for extraction. The embedding process hides a secret message within a cover media (cover image), and the result of embedding process is stego image. The main issue is that the secret message will not be unnoticed if a third party tries to intercept the cover media (cover image). The extraction process is simply because it is the inverse of the embedding process, where the secret message is revealed at the end [8]. To evaluate the quality of image, stego image and cover image are compared. This requires a measure of stego-image quality, commonly used measures are Mean Square Error (MSE), and Peak Signal-to-Noise Ratio (PSNR). Mean Square Error (MSE) is used to quantify the difference between the initial (cover) and the distorted or noisy (Stego) image [8, 9].

$$MSE = \frac{1}{MN} \sum_{x=1}^M \sum_{y=1}^N (S_{xy} - C_{xy}) \dots\dots\dots (1)$$

Where X and Y are the image coordinates, M and N are the number of rows and columns in the input images, respectively. Sxy is the generated stego-image and Cxy is the cover image [7].

Peak Signal-to-Noise Ratio (PSNR) is used to measure image distortion due to embedding and it is measured in decibels (dB) [9, 10].

$$PSNR=10\log_{10} (C_{2max}/MSE)\dots\dots\dots(2)$$

Where, Cmax holds the maximum value in the image that is 255 and MSE is the mean square error which is determined by equation (1).

The steganography approaches can be divided into three types [11]:

- 1) *Pure Steganography*; it is a technique simply uses the steganography approach only without combining other methods. It is working on hiding information within cover carrier.
- 2) *Secret Key Steganography*; it uses the combination of the secret key cryptography technique and the steganography approach. The idea of this type is to encrypt the secret message by secret key technique and then hide the encrypted data within cover carrier.
- 3) *Public Key Steganography*; it is the combination of the public key cryptography approach and the steganography approach. The idea of this type is to encrypt the secret data using the public key approach and then hide the encrypted data within cover carrier.

The Difference between Cryptography and Steganography [8, 11]:

- Cryptography prevents unauthorized party from discovering the content of communication but Steganography prevents discovery of the existence of communication (i.e., Cryptography makes data gibberish and known the message passing while Steganography tends to conceal presence of hidden data and unknown the message passing).
- Cryptography alters the structure of secret message while Steganography does not alter the structure of secret message.
- Cryptography is more common technology than Steganography technology.

- The most algorithms of Cryptography are well known, but the algorithms of Steganography are still being developed by certain formats.
- In Cryptography, the strong algorithm depends on the key size, the more key size; the more expensive computing power is required to decrypt ciphertext. In Steganography, once the hidden message is detected, the message is become known.
- Cryptography can provide all security objectives by implementing the public and private key(s) with hash functions or authentication codes or digital signatures. Steganography cannot provide most of security objectives (Integrity, authenticity, non-repudiation) by itself without using the cryptographic techniques. However, it provides confidentiality by itself because mostly, the concerning person knows that the message is hidden in what kind of medium [12].

In this paper, the secret Key steganography approach is used to improve security by using modified AES and method in [1] which includes PVD\_MPK and MSLDIP-MPK methods to encrypt and hide the message in cover image. Therefore, if an attacker doubts about the stego image and tries to detect the message from the stego image, he would still require the key to decrypt the encrypted message.

## II. RELATED WORK

There are many methods have been used to provide data security whether by using encryption, steganography or combination between them. Advance encryption standard (AES) method, it is also known as Rijndael, is a symmetric-key block cipher [12]. Unlike DES method, AES method is a non-Feistel cipher that encrypts and decrypts a data block of 128 bits. It uses 10, 12, or 14 rounds. The key size, which can be 128, 192, or 256 bits, and the number of rounds depend on the key size because it allows the secret key to be expanded to produce sub key for each round. In AES method, the input and output sequences have the same length [12, 13]. According to AES method, substitution byte, shift rows, mixing column and key adding steps are implemented in every encryption round to encrypt the message, but the Mixing Column step doesn't include in the last round. In the decryption, the four steps are implemented in the reverse way. Also, the inverse of mixing column step doesn't include in the last round of the decryption. The pseudo code of AES is as follows [13]:

```

1. InitialRound(State, RoundKey)
{
    AddRoundKey (State, RoundKey);
}
2. Rounds (State, RoundKey)
{
    SubBytes (State);
    ShiftRows(State);
    MixColumn(State);
    AddRoundKey(State, RoundKey);
}
3. FinalRound(State, RoundKey)
{
    SubBytes (State) ;
    ShiftRows(State) ;
    AddRoundKey(State, RoundKey);
}

```

Where SubBytes is a non-linear substitution step where each byte is replaced with another according to a lookup table after each byte is interpreted as two hexadecimal digits. ShiftRows is a transposition step where it performs a circular rotates on each row of 0, 1, 2 & 3 places for respective rows. MixColumns is a mixing operation which operates on the columns of the state. It transforms each column of the state to a new column. AddRoundKey is precedes one column at a time. AddRoundKey adds a round key word with each state column matrix using bitwise XOR operation [14]. The advantages of using AES algorithm are; it is more secure, support larger key sizes than DES, faster in both hardware and software, reasonable cost, and its main characteristics flexibility and simplicity [15]. In [16], a method has been proposed based on spatial domain instead of using LSB1 (First Least Significant Bit) of the cover image for embedding the message bits, LSB-3 (Third Least Significant Bit) has been used to hold the message bits and LSB-1, LSB-2 may also be modified. According to this method, it is found that the LSB-1 method has more PSNR (Peak Signal to Noise Ratio) values which mean that the image's quality using LSB-1 method is better than that LSB-3 method. Also, the capacity is still the same. So, the basic LSB1 method has better results than the proposed LSB-3 method. In [1], a steganography algorithm has been proposed by combining PVD-

MPK and MSLDIP-MPK methods. The idea of this proposed algorithm is that using the digits of MPK encoding instead of bits of secret message in hidden step to increase the capacity and quality because MPK encoding represents each character by two digits instead of 8 bits. As an example, the letter "a" will be represented as 2 1, this mean that the letter "a" can be typed by pressing the key no.# (2) in the keypad only one time. Firstly, the cover image is divided into non-overlapping blocks, where each block consists of two consecutive pixels and each byte in the message is converted to Mobile Phone Keypad (MPK) format, then the difference value  $d_i$  for each block of two consecutive pixels  $P_i, P_{i+1}$  is calculated. If the difference value is larger than 19, this means that the two pixels are located in high level (edge area) and MSLDIPMPK method is used to hid data to increase capacity. Therefore, each last digit in two pixels of block is replaced by the digit value of the secret message. On the contrary, if the difference value is smaller than 19, this means that the two pixels are located in low level (smooth area) and PVD\_MPK method is used to hide data to increase image quality. Then, find the optimum range  $R_i$  for the difference  $d_i$  where  $R_i = [l_i, u_i]$ , and  $l_i$  and  $u_i$  are the lower and upper bound of each range of the range table and  $l_i \leq d_i \leq u_i$  because this method uses the ranges (0-4), (5-9), (10-14), (15-19), (20-24), (25-29) etc. to calculate the new difference and make it nearly as possible to the original difference. The new difference is calculated by equation (3):

$d'_i = l_i + b/2 \dots \dots \dots (3)$  Where  $d'_i$  is the new difference,  $l_i$  is the lower bound of the range, and  $b$  is the digit that needs to be hidden. Replace the old difference with the new difference by using Equation (4):

$$(p'_i, p'_{i+1}) = \begin{cases} (p_i + \lfloor \frac{m}{2} \rfloor, p_{i+1} - \lfloor \frac{m}{2} \rfloor), & \text{if } p_i \geq p_{i+1} \text{ and } d'_i > d_i; \\ (p_i - \lfloor \frac{m}{2} \rfloor, p_{i+1} + \lfloor \frac{m}{2} \rfloor), & \text{if } p_i < p_{i+1} \text{ and } d'_i > d_i; \\ (p_i - \lfloor \frac{m}{2} \rfloor, p_{i+1} + \lfloor \frac{m}{2} \rfloor), & \text{if } p_i \geq p_{i+1} \text{ and } d'_i \leq d_i; \\ (p_i + \lfloor \frac{m}{2} \rfloor, p_{i+1} - \lfloor \frac{m}{2} \rfloor), & \text{if } p_i < p_{i+1} \text{ and } d'_i \leq d_i. \end{cases} \dots (4)$$

This equation adjusts the value of  $p_i$  and  $p_{i+1}$  to modify the difference. Where,  $m$  is the difference between the original and the new differences (i.e.,  $m = |d_i - d'_i|$ ),  $p_i$  and  $p_{i+1}$  are the first and second pixel in a block before embedding, and  $p'_i$  and  $p'_{i+1}$  are the first

and second pixel in a block after embedding. Then, Compute the remainder value by:

$$\text{Rem} = b \text{ mod } 2 \dots \dots \dots (5)$$

If Rem equal 1, make  $p_i$  odd, else make  $p_i$  even. The experimental results showed that the PSNR and Maximum Hidden Capacity (MHC) of this technique are high comparing to the existed techniques. In [17], a proposed technique has been introduced. The results of the proposed technique prove that it is more secure than other techniques against statistical attacks which are commonly used in steganalysis. This is because it ranked images in a user's library based on their suitability as cover objects for some data. The process is repeated for each image in a user's image library. Each bit of the encrypted data is compared to the least significant bit of the pixel bytes in an image. The comparisons are made starting from the first byte until the last byte in the image that permits all data to be hidden in that image. By matching data to an image, there is less chance of an attacker being able to use steganalysis to recover the data. The image is then given a rank based on the percentage of least significant bits that match the encrypted data bits. Before hiding the data in an image, the data is first encrypted using the RSA public key algorithm, and then the encrypted data is hidden by using LSB. In [18], a new method has been proposed for hiding any encrypted secret message inside a cover file. To embed a secret message file in the cover file, they have used two distinct methods; (1) they encrypt the secret message file using simple bit shifting and XOR operation in the secret message file. (2) The encrypted secret message is embedded in the cover file in alternate byte instead of changing the LSB of the cover file bytes. They change LSB and LSB+3 bits of the cover file bytes. Their method could be most appropriate for hiding any file in any standard or nonstandard cover file such as word, excel, .ppt, .exe, image, audio, video files. Also, this method may be used for sending some secret key to someone over mail as the intruder may not be able to unhide and decrypt the secret message.

### III. PROPOSED METHOD

The main objective of the proposed method is to introduce more secure communication by merging cryptography and steganography techniques to make it more difficult for a steganalyst to retrieve the plaintext of a secret message from a stego-object. The proposed method is divided into two parts. In the first

part, the AES algorithm will be modified to be suitable for steganography method and it is called AES\_MPK algorithm. In the second part, the AES\_MPK algorithm will be merged with steganography algorithm to hide the encrypted data in image where a message being sent is concealed [1]. Therefore, two levels of security have been applied.

A. *The Modified AES (AES\_MPK) Algorithm*

According to the modified AES\_MPK algorithm, four types of transformations are used like AES; substitution (SubBytes), permutation (ShiftRows), MixColumns, and key-adding, to provide security [13]. Because of the AES is based on the Rijndael cipher, it performs four types of transformation based on the operations in finite field GF (2<sup>8</sup>). Several operations are defined at byte level, and used with bytes representing as elements in the finite field or Galois field GF (2<sup>8</sup>). Then, it represents the input and the output in form of hexadecimal digits (i.e. two hexadecimal digits for each byte) [13]. Therefore, AES algorithm is modified to make the input and the output in the form of MPK digits because the PVD\_MPK and MSLDIP-MPK methods use the MPK digits for hiding the data. This modified AES algorithm called AES\_MPK algorithm.

The basic idea of AES\_MPK algorithm is to convert each block (i.e. 16 byte) of data to MPK digits (i.e. two digits for each byte). Then each block is divided into two states; the first state represents the first digits of 16 byte and second state represents the second digits of 16 byte. Two states are filtered by replacing each digit 8 or 9 with 7 to make state suitable for operations of GF (2<sup>3</sup>) that are applied at each stage of encryption but save them in their location in each state. Finally, apply operations of GF (2<sup>3</sup>) in MPK digits in the four types of transformation expect 8 and 9 digits which will be replaced with 7 and keep in their location. This is because the elements of GF (2<sup>3</sup>) are 0, 1, 2, 3, 4, 5, 6, and 7. The four types of transformation are: 1) *Substitution*: The first transformation is SubBytes. To substitute a byte, this requires two steps; first apply multiplicative inverse in GF (2<sup>3</sup>) on each element in two states. Second, apply an affine transformation over GF (2<sup>3</sup>) on each element in the two states by representing each element as 3 bits (b<sub>0</sub>, b<sub>1</sub>, b<sub>2</sub>). AES\_MPK depicts this transformation in matrix form as follows:

$$\begin{bmatrix} c_0 \\ c_1 \\ c_2 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \\ 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} b_0 \\ b_1 \\ b_2 \end{bmatrix} + \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix}$$

Where (b<sub>0</sub>, b<sub>1</sub>, b<sub>2</sub>) are bits of element before affine transformation and (c<sub>0</sub>, c<sub>1</sub>, c<sub>2</sub>) are bits of element after affine transformation.

The SubBytes stage provides confusion (i.e., it obscures the relationship between the plaintext and the ciphertext) [5]. In decryption, InvSubBytes is performed by applying affine transformation first, and then multiplicative inverse in GF (2<sup>3</sup>) is applied. The affine transformation in decryption will be as follows:

$$\begin{bmatrix} b_0 \\ b_1 \\ b_2 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix} \begin{bmatrix} c_0 \\ c_1 \\ c_2 \end{bmatrix} + \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix}$$

2) *Permutation*: Another transformation found in a round is called ShiftRows, which permutes the bytes. It performs a circular rotate on each row of 0, 1, 2 & 3 places for respective rows. This shift moves an individual byte from one column to another, and ensures that the 4 bytes of one column are spread out to four different columns. The ShiftRows stage provides diffusion of column values between columns. The rows shift to left as follows:

- Row0: No shift
- Row1: 1-byte shift
- Row2: 2-byte shift
- Row3: 3-byte shift

In decryption, InvShiftRows is performed where the circular will be shifted in the opposite direction for each row.

3) *MixColumns*: This stage is used to mix bytes by using matrix multiplication. It operates at the column level where each column of the state is transformed to a new column by multiplying it in constant matrix. This stage changes the bits inside a byte, based on the bits inside the neighboring bytes. We need to mix bytes to provide diffusion at the bit level.

In decryption, InvMixColumns is performed by the inverse of constant matrix in multiplication.

$$C = \begin{pmatrix} 2 & 3 & 1 & 1 \\ 1 & 2 & 3 & 1 \\ 1 & 1 & 2 & 3 \\ 3 & 1 & 1 & 2 \end{pmatrix} \xrightarrow{\text{Inverse}} C^{-1} = \begin{pmatrix} 5 & 0 & 6 & 2 \\ 2 & 5 & 0 & 6 \\ 6 & 2 & 5 & 0 \\ 0 & 6 & 2 & 5 \end{pmatrix}$$

Where C is constant matrix and C-1 is inverse constant matrix in GF (2<sup>3</sup>)

4) *AddRoundKey*: the AddRoundKey stage is a simple bitwise XOR of the current block (two states) with two portions of the expanded key. Note; this is the only step which makes use of the key. This step makes the cipher as a series of XOR with expanded keys then scramble/permute block repeated. This provides efficient and highly secure.

Note that: AES has defined three versions, with 10, 12, and 14 rounds. Each version uses a different cipher key size. We work in version with 10 rounds and key size is 16 byte.

The pseudo code of the modified AES\_MPK algorithm is as follows:

**AES\_MPK Algorithm**

**Input:** Secret Message SM, Cipher Key K.

**Output:** Cipher Message CM.

**Steps:**

1. Make key expansion of K that produces two lists of all sub keys.
2. Partition SM to blocks (B<sub>1</sub>, B<sub>2</sub>, B<sub>3</sub> ... B<sub>n</sub>) each block consists of 16 byte.
3. **for** each B<sub>i</sub> block **do**
4. Convert each byte to MPK digits (two digits for each byte).
5. Divide B<sub>i</sub> to two state arrays (4\*4).
6. Filter two states.
7. Make pre round AddRoundKey which is a simple bitwise XOR of the current two states with two sub keys
8. **repeat**
9. Apply the four transformations (SubBytes, ShiftRows, MixColumns, and AddRoundKey) in two states.
10. **until** nine round.
11. At final round implements SubBytes, ShiftRows, and AddRoundKey but MixColumns is deleted.
12. Return the digits 9 and 8 in their place in each state.
13. Mix two states to be one block.
14. Convert block to characters by using MPK decoding (i.e. two digits represent character). The result represents cipher block
15. **end**
16. Concatenate the currently cipher block with the previous cipher blocks to collect CM.

V. CONCLUSIONS AND FUTURE WORK

In this paper, a new secure communication model has been presented that combines cryptography and steganography techniques to provide two layer of security, so the steganalyst can't reach to plaintext without knowing the secret key to decrypt the ciphertext. Firstly the secret data has been encrypted by using the AES\_MPK then the encrypted data has

been hidden in gray image by using PVD-MPK and MSLDIPMPK methods. Due to this combination, the secret data can transmit over open channel because the cipher text does not look meaningless but its presence is concealed by using steganography for hiding cipher text in the images. Experimental results showed that our proposed model can be used to hide much more information than that other existed methods and the visual quality of the stego image is also improved, in addition to it is effective for secret data communication.

In the future work, we are looking forward to try applying the proposed method on audio and video. Also, we are looking forward to enhance the proposed method to make the capacity higher than it while keeping the same PSNR or higher.

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