SESPHR:AMETHODOLOGYFORSECURESHARINGOFPERS ONALHEALTHRECORDSINTHECLOUD

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ABSTRACT

Thebroaduseofcloud-

basedservices in the health care industry has made it pos sible for various participating entities of the e-

Healthsystemstoexchange personal health records (PHRs) at a lowcost and with ease. However, putting the privatehealth data on cloud servers leaves it open to theftordisclosure, necessitating the creation of proced uresthatprotectthePHRs'privacy.Consequently,wes uggestatechniquenamedSeSPHR for PHR cloud sharing that is secure. TheSeSPHRsystemmakessurethatPHRsarecontroll ed from а patientcentric perspective and maintains the maintaining the PHRs' privacy. Patie encrypted PHRs on unreliable nts keep cloudserversandonlygivecertainpeopleaccess.onva riousPHR sections, to distinct categories of users. a and re-encryption semisetup server, a trustedproxy(SRS)isusedtocreatethereencryptionkeysandtocreatethepublic/privatekeypai

rings.Additionally,theprocessissafe.It

implementsaforwardandbackwardaccesscontroland protects against insider risks. Additionally, weofficiallyevaluateandconfirmtheuseoftheSeSPH Rapproachusinghigh-levelpetrinets(HLPN).

Evaluation of performance in relation totimeConsumptionsuggeststhattheSeSPHRapproa chmaybeusedforsafelytransferringPHRstothecloud

I. INTRODUCTION

Inordertoprovideubiquitousandon-

demandavailability of different resources in the form

ofhardware,software,infrastructure,andstorage,CL OUD computing has developed as a significantcomputing paradigm [1, 2]. As a result, the cloudcomputing paradigm helps enterprises by relievingthem of the time- consuming task of developinginfrastructureandencouragingthemtorel yonoutside Information Technology (IT) services [3].Additionally, the cloud computing architecture hasshowntremendouspromiseforimprovingcoordin ationamongmanyhealthcarestakeholders andforguaranteeingscalabilityandongoingavailabili tyofhealthinformation[4,5].Additionally,thecloudc omputingconnectsanumberofsignificanthealthcare domains,includingpatients,hospitalstaff,includingp hysiciansandnurses,pharmacists,andclinicallaborat orystaff,aswellasinsurancecompaniesandservicepro viders[6].Asaconsequence,acollaborative and costeffective health ecosystemwhere patients may easily establish and maintaintheir Personal Health Records (PHRs) develops as result of the integration of the aforementionedorganisations[7].ThePHRsoftenincl udedatalike:

(a) demographics, (b) medical history,includingdiagnoses, allergies, surgeries, andtreatments,

(c)laboratoryresults,(d)informationonhealthinsuran ceclaims,and(e)patient-

onlynotesregardingspecificsignificantobservedheal thissues [8].

Moretechnically,PHRsarecontrolledthroughInterne t-based technologies, allowing individualsto manage their health information as permanentrecords that can be accessed by those who need it[9]. As a result, PHRs make it possible for

peopletosuccessfullycommunicatewithmedicalprof essionals in order to describe their symptoms,ask for guidance, and maintain their health recordsforproper diagnosis and treatment.

Despitethebenefitsofthescalable,adaptable,afforda ble,andwidespreadservicesprovidedby the cloud, a number of issues linked to the privacyofhealthdataalsocomeup.Theuseofthecloudt odistributeandstorePHRsisacrucialfactorinpatients' concerns about the confidentiality of suchrecords [10]. Private health information stored oncloud servers run by third parties is vulnerable tointrusion. Particularly jeopardised is the privacy ofPHRskeptinpubliccloudsrunbyfor-

profitserviceproviders [11]. The PHRs' privacy may be underdanger in a number of ways, including theft, loss, and leaking [12]. Because of the male volent action sofotherentities, the PHR sincloudstorage, intransit from the patient to the cloud, or from the cloud to any other user, may be vulnerable to unauth orisedaccess.Additionally,thereareoccasional threats made against the data by realinsiders [13]. For instance, due of the nefariousactions of other organisations, the PHRs in cloudstorage, in transit from the patient to the cloud. orfromthecloudtoanyotheruser, maybevulnerableto illegal access [10]. People who work for the cloud service provider may act maliciously. Theepisode

employee took home without authority theprivatehealthinformationofover26.5millionpeop leisawell-known illustrationof that[14]. TheHealthInsurancePortabilityandAccountability

in which a U.S. Department of VeteransAffairs

Act (HIPAA) requires that patients'consent and the terms of use and disclosure befollowed in order to maintain the integ rity and confidentiality of electronic health informatio byhealthcareproviders[15].Additionally,thePHRs shouldbeencryptedwhenbeingkeptonthird-party cloud storage so that neither the cloudserverprovidersnorunauthorisedpartiesmayac cessthePHRs.ThePHRsshouldonlybeaccessibletoe ntitiesorpeoplewhohavethe"right-to-

know"privilege.Topreventunauthorisedalterationso rabuseofdatawhenitistransferredtotheotherstakehol dersinthehealthcloudenvironment, the mechanism for granting access toPHRsshouldbemanagedbythepatientsthemselves.

TheprivacyofPHRskeptoncloudservershasbeenprot ectedinavarietyofways.Confidentiality,integrity, authenticity,accountability,

and audit trial are ensured by privacypreservingmethods. While integrity concerns with preservingthe originality of the data, whether in transit or incloud storage, confidentiality guarantees that thehealthinformationiscompletelyhiddenfromunaut horisedparties[14].

Accountability refers to the need that data accessregulations follow the established protocols, whileauthenticity ensuresthat the health data isonlyaccessible by authorised parties. The term "audittrial" refers to the process of observing how healthdata is used even after access to it has been given[6].

We provide a way for managing the PHR accesscontrolsystemthatiscontrolledbypatients

themselves, dubbed Secure Sharing of PHRs in theCloud (SeSPHR).

approach limits unauthorised The users to protectthePHRs'confidentiality.Inthesuggestedappr oach, there are typically two categories of PHR users: (a) patients or PHR owners; and (b) users of PHRs who are not owners. such as patients' familymembers or friends, physicians, health insurancecompanyrepresentatives, pharmacists, and researchers.

By selectively providing people access to certainPHR sections, patients who are the PHRs' ownersareallowedtouploadencryptedPHRstotheclo ud.Depending on their job, each member of the groupof users of the latter kind is given access to thePHRs to a certain degree by the PHR owners. ThePHR owner defines the degrees of access given todifferent user groups in the Access Control List(ACL).

For instance, the owner of the PHRs may providecomplete access to the patient's family members oracquaintances. Similar this, to insurance companypersonnel would only be allowed the to see PHRsectionsthatincludeinformationconcerninghea lthinsuranceclaims, with access to other personal medi calinformation, such the patient's medical history, bein gblockedfortheseusers.

The SeSPHR methodology avoids the overhead bydelegatingtheSRSforsettingupthepublic/privatek eypairsandproducingthedecryptionkeysforthe authorised users only. In contrast to the approachproposedin[10], which suggests that the PH Rownersmanagemultiplekeys, this approach avoid so verheadbyproposingthatthePHRownersproduce the decryption keys. This ultimately leadstooverheadatthePHRowner'send.TheSetupand Reencryption Server (SRS), a semi-trusted server, is used as the proxy since the approach view sclo udserversasanuntrustedentity.FortheSRStogenerate the re-encryption keys for safe sharing ofPHRsacrossusers, aproxyreencryption-

basedtechniqueisutilised.PatientsorPHRownersenc rypt the PHRs, and only authorised users withkeys provided by the SRS may decode the PHRs.Additionally, the users are given access to the PH Rs'particularsectionsthatthePHRownerdeemstobec rucial.Theproposedmethodissecurecomparedtoprev iousconstructssincethePHRdataisneversentfromthe SRSintheproposedframework.Instead,itistheSRS's dutytomaintainthe keys, with PHR owners handling

encryptiontasks,andrequestingusershandlingdecryp tiontasks,providedtheyhaveaccesstovaliddecryptio nkeys.

Theforwardandbackwardaccesscontrolsarelikewise enforced by the suggested method. Thekeysareobtainedbythenewlyjoiningmembersof acertainusergroupfromtheSRS.Onlytheowner'skey s areusedto encrypttheshared data. After receiving the PHR owner's consent, newlyjoining members are given access to the data. The corresponding keys for a departing user are also destroyed, and that user is also removed from

theACL.Anyunauthorisedaccessattemptsmadeafter the user has left are denied access to the PHR dueto the deletion of the user keys and removal from the

ACL.WealsousedHighLevelPetriNets(HLPN)andt heZlanguagetodoaformalexaminationofthesuggest eddesign.

TheHLPNisusedtobothimitatethesystemandtoprovi de the mathematical characteristics that arelaterutilisedtoanalysethebehaviourofthesystem. TheZ3solverandtheSatisfiabilityModuloTheories (SMT-Lib) are Library used to carry outtheverification. Tocarryouttheworkofverificatio SMT, the petri net model n using the isfirsttranslated into the SMT together with the specifie properties. and then the Z3 solver is d usedtocheckwhethertheproperties aretrueorfalse.



Fig-

1:ArchitectureoftheproposedSeSPHRmethodology jespublication.com Page 701 The following are the main contributions of thesuggestedwork:

1. SeSPHR, a technique we offer, enables patients control the sharing of their own PHRs in the cloud.

2. TomaintainPHRsecrecy, the SeSPHR approachus esproxyre-encryption and El-Gamalencryption.

3. Based on the access level established in the ACL for various user groups, the approach enables PHRowners to selectively provide users access to usersoverthesections of PHRs.

4. To provide access control and to produce thereencryption keys for various user groups, a semi-

trustedproxynamedSRSisimplemented,removing the burden of key management from thePHR owner'send.

5. The suggested technique also uses forward andbackwardaccesscontrol.

6. Thesuggested approachis formally analysed and ve rified to ensure that it operates in accordance with there quirements.

II. RELATEDWORK

The existing works that are related to the proposedwork are presented in this section. By sending

thePersonallyIdentifiableInformationindividually,t heauthorsin[28]developedapublickeyencryptionbasedtechniquetomaintaintheanonymityandunlinka in a semitrusted cloud (PII). The Cloud ServiceProvider(CSP)savesthe healthrecordandthelocation of the file (index), and later encrypts themusingsymmetrickey encryption.

ThepatientsencryptthePHRsbythepatientsusingthep ublickeyoftheCSP,andtheCSPdecryptstherecord using the private key. By associating thelocationandthemasterkey,theadministrativecont rol of the patient on the PHRs is kept in place.The approach's drawback is that it enables the CSPtodecryptPHRs,whichmaythenbeusedmaliciou sly.TheSRS,ontheotherhand,isasemi-

trustedauthoritythatdecryptstheciphertextcreatedby thePHRownerandprovideskeystotheusers thatrequestaccesstothePHRs.

Inamulti-usercloudsetting, Chenetaltechnique.'s [12] uses the SKE and the Lagrange Multiplier todynamically exercise access control on PHRs. Theapproach'sprimarycharacteristicsincludeautom aticuserrevocation. Apartial orderlink between the users is kept in order to get around thechallengesofotherkeymanagement. However, the systemrequiresthePHRowners'onlinepresenceinord ertogiveorcancelaccess.Oursuggestedtechnique does not need the PHR owners to be nline in order to offer the access over PHRs, incontrast to the plan described in [12]. Instead. thesemitrustedauthoritychoosestheusers'accessrightsand,u ponsuccessfulauthorisation, decides

bilityofhealthinformation ISSN:0377-9254 the re-encryption keys for the users making therequest.

Toprovidepatient-

centricaccesscontrol,theauthorsin[29]developedaD igitalRightManagement (DRM)-based solution. The

writersusedContentKeyEncryption(CKE)toencrypt thedata,andonlyuserswithvalidlicencesarealloweda ccess.Thefirstproxyre-

encryptiontechniquewasputoutin[33].Unlikeourpol icy,whichisbasedonkeysandhasnoeffectonthesizeof theciphertext,the policy in [33] is based on ciphertext, and thesize of the ciphertext rises linearly with multi-useusage. This is because the [33] needs a step that ismissing from our methodology—re- encryption. Lietal.[14]offeramethodforsharingPHRsinmultiownersettingsthatareseparatedintoseveraldomainst

hatusesattribute-basedencryption(ABE).

The technique initially presented in [33] serves asthefoundationforthesuggestedmethodology.After a given user's access has been revoked, themethod re-encrypts the PHRs using the proxy reencryptionmechanism(s).Themethodsuccessfullyre ducesthecomplexityandexpenseofkeymanagement whilealsoimprovingthephenomena of on-demand user revocation.

Despitebeingscalable,themethodisunabletohandlesi tuations when granting access permissions basedonusers' identities is necessary. To guarantee user responsibility, Xhafa et al. [30]also applied Ciphertext Policy ABE (CPABE). Inaddition to preserving user privacy, the suggested method has the ability to track down users

whomisbehaveandillegitimatelysharetheirdecrypti onkeys with other users.

It presents a method for ensuring both the secrecyand fine-grained access to the healthcare data thathas been contracted out to cloud servers. By usingproxy re-encryption, Key Policy ABE (KP-ABE), and lazy re-encryption, the expensive of duties reencryptingdatafiles, updatingsecretkeys, and prevent ing users whose access has been revokedfrom learning the of the contents data are handled.There-

encryptionofdatafilesandsubsequentstorageinthecl oudenvironmentareresponsibilitiesassignedtothecl oudservers.However, the data owner is also expected in theproposed framework to be a reliable authority whocontrols thekeys forseveralowners andusers.

Therefore, managings everal keys for various attribute s for many owners would be inefficient atthe PHR owners' end. Because the functions of key cre ation and key distribution to various user types are carried out by the semi- trusted authority, our technique eliminates over head. In order to provide fine-grained access control, the authors in [31] and

[32]alsoemployedproxyre-encryption-

jespublication based techniques. The system we provide Ralgev703

owners to encrypt PHRs before putting the minthe clou

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and adds a semi-trusted authority that reencryptsthe ciphertext without knowing what's within

thePHRs.ThePHRscanonlybedecryptedbyauthorise duserswhopossessdecryptionkeysissuedbythesemitrustedauthority.

III. PROPOSEDMETHODOLOGY

The recommended method makes use of proxy reencryptiontoprovidePHRconfidentialityandexchan gesecurityacrosspublicclouds.Thearchitecture of the suggested SeSPHR technique isshownin Fig. 1.

Persons the recommended method for exchangingPHRs in a cloud environment involves the SetupandReencryptionServer(SRS),thecloud,andtheusers.Ano verviewofeachoftheentitiesisprovidedbelow.

The cloud The strategy advises PHR owners to save their r data in the cloud so they may subsequently safely share it with other users. Users assume that the cloud is an unreliable source when they upload or download PHRs to or from cloud servers. No changes to the cloud are necessary since both types of fusers are the only one sthat upload and download PHRs in the way stated.

Setting up and installing the SRS: Every systemuser'spublic/privatekeypairsmustbegenerate dbythe SRS, a semi-trusted server. The SRS furthergenerates the re-encryption keys in order to safelytransferPHRamongseveralusergroups.TheSR isregardedasasemitrustedentityintherecommended method.Inlightofthis,wedrawtheconclusion that it is honest and generally followsthelaw,butodd.TheSRSmonitorsthe keys,butitnevergetsPHRinformation.Operationsfor encryptionanddecryptionarecompletedattheendpoi ntsoftheusers.TheSRSofferskeymanagement in addition to access control for theshareddata.

Dueofthepubliccloud'sunreliability,theSRSisastand aloneserverthatcannotbeinstalledthere.TheSRS may be managed by a group of institutions orby a respectable third-party organisation forthebenefit of the patients. It could also be maintainedby a group of connected patients. However, SRSmaintainedbyhospitalsoragroupofpatientsmigh tinspirehighertrustduetotheinvolvementofmedical specialists and/or the patients' self-controlover SRS.

Users:Patients(ownersofPHRswhowishtosecurelys haretheirPHRswithothers)andpatients'familymemb ersorfriends,doctors,representativesof health insurance companies, pharmacists, andresearchers are the two main groups of users of thesystem. Friends and relatives are classed as privatedomainusersundertheSeSPHRmethodology, whereas all other users are categorised as publicdomainusers.

PHR owners may provide users access to PHRs inthepublicandprivatedomainstovaryingdegrees.

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Userswhocomeundertheprivatedomain, for example complete may have access to the PHR, but those who fall under the public domain, such a sphysicians, scientists, and pharmacists, would onlyhave access to a limited number of PHR parts. Theaforementionedusersmayalsobegrantedfullacce ss to the PHRs if the PHR owner determines itis required. In other words. the SeSPHR approachallows patients to impose precise access controlover PHRs.

Every system user must register with the SRS inorder to access the SRS's services. As a doctor,researcher,orpharmacist,forexample,theregi stration procedure is dependent on the user'sresponsibilities.

HRPartitioning

The four sections listed below are logical divisions of the PHR:

Personalinformation, health-

related information, insurance information, and infor mation on prescription drugs;

Itiscrucialtonotethattheaforementioneddivisionis flexible. The PHR may be divided into fewer ormore divisions at the user's discretion. The PHRsare represented in a number of formats, includingXML,andmaybesimplyseparatedintopiece s.ThePHR owner also has the option of giving manypartitions the same level of access control. SomePHRpartitionsmayincludeuserrestrictions, meaning that a particular user may not have fullaccess to thehealth data.

For instance, a pharmacist may not have access topersonal or medical information, but they may begivenprescriptionandinsurance-

related information. Full access to the PHR may also begiven to family members and friends. A researchercould only need access to the patient's medic alrecords once the personal data has been deleted. The PHR owner grants the SRS access rights to each of the various PHR partitions when data is uploaded to the cloud.

The Proposed Methodology's Approach Functions

The suggested SeSPHR technique consists of thefollowing steps: setup, key creation, encryption, and decryption. The parts that follow go througheachaction:

SETUP

The offered approaches work well with the G1 andG2 groups with the prime order q. G1 G1G1 andG2 are bilinearly mapped to form G2. A randomnumbergeneratorwheregG1hasgasaparamet er.Z is used as a second random number generatorusingtheformulaZ= e(g, g) G2.

KeyGeneration

Public/privatekeypairsarecreatedbytheSRSforthes etof authorised users.

ENCRYPTION

Imagine that patient P is required to upload theirPHR to the cloud. The PHR partitions that the userhasallocatedtothedifferentaccesslevelgroupsar erepresented by a random number or numbers thatare generated by the patient client application. Inour case, we take into consideration that the accesslevels for each of the four partitions specified inSection 3.2 vary. As a consequence, four randomvariables are created in our example: r1, r2, r3, andr4 (Zq). The variable ri is used to encrypt the i-thpartition of the PHR. Each partition is encryptedseparately by the client programme. Thanks to

theXMLstructure,theapplicationcanquicklyperform encryptionanddecryptiononthe

PHR'slogicalpartitions. The partitions stated above in the PHRareencryptedasseenbelow.

The quantity of PHR partitions; the titles of eachpartition, such as "Personal Health Record" (PHR), "Medical Record," "Insurance Information," and "Pr escription Information; (any role may be granted access to more than one partition, such as doctors may be granted access to medical information).

• The first close relative or acquaintance toprovide access • If there is any default access fornew members

DECRYPTION

Let's assume user U requests access to the patientP'sprovidedencryptedPHR(C).UserUdownl oadsthe C directly from the cloud after completing thecloudauthenticationprocess.TheuserUrequests

the SRS to determine and deliver the correct Rparameters required for decryption at that point. By lo okingattheaskinguser'sACL,theSRSdetermines if the PHR owner has granted access to he partition for which the user has requested R.BasedontheaccessrightsspecifiedintheACL,theS RSwillcreateandprovidethenecessaryparameters to the requesting user. We shall showhowRisproducedforeachdivisioninthetextthat followsinordertoprovideacomprehensiveexplanatio noftheprocedure. Therefore, we assume that user U has complete access to all partitions. The SRS computes R and sends it to the user Utogetherwiththere-encryptionkey.

VI.CONCLUSION

We suggested a mechanism for transmitting andstoring PHRs in the cloud securely to authorizedparties.ThetechniqueupholdsapatientcentricaccesscontroltovariousPHRsubsystemsbase dontheaccessgrantedbythepatients,whilemaintainin gtheprivacyof thePHRs.

Weputinplaceaformoffine-grainedaccessrestriction so that not even authorized users of thesystem could access restricted areas of the PHR.Onlyauthoriseduserswithlegitimatere-

encryption keys supplied by a semi trusted proxyareabletodecryptPHRs,whicharestoredencryp ted by PHR owners in the cloud. The tractortrailer proxy's job is to create and maintain publicandprivatekeypairsforthesystem's users. Themethodologyalsomanagesforwardandbackward identitymanagementforleavingandnewlyjoininguse rs,correspondingly,inadditiontomaintainingconfide ntialityandguaranteeingpatient-

centricaccesscontrolforPHRs.Additionally, we officially assessed and validatedtheSeSPHRmethodology'soperationusingt heHLPN,SMT-

Lib,andZ3solver.Thetimeittooktogeneratekeys,thea ctivitiesinvolvedinbothencryptionanddecryptioninc ludingtimelydeliverywasalltakenintoaccountwhene valuatingperformance.Theoutcomesoftheexperime nt show that the SeSPHR approach maybeusedtosafelyexchangePHRsinacloudcontext.

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